

Logical foundations of categorization theory

Handout 1

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1 Preliminaries

Definition 1 (Partial order). *Let P be a set. A partial order on P is a binary relation \leq on P such that, for all $x, y, z \in P$,*

- (a) $x \leq x$ (reflexive);
- (b) if $x \leq y$ and $y \leq x$ then $x = y$ (antisymmetric);
- (c) if $x \leq y$ and $y \leq z$ then $x \leq z$ (transitive).

If \leq is a partial order on P , we call (P, \leq) a partially ordered set (poset).

Definition 2 (Closure operator). *If (P, \leq) is a partially ordered set, a map $c : P \rightarrow P$ is a closure operator on (P, \leq) , if for all $x, y \in P$,*

- (a) $x \leq c(x)$ (inflationary);
- (b) if $x \leq y$ then $c(x) \leq c(y)$ (monotone);
- (c) $c(c(x)) = c(x)$ (idempotent).

If c is a closure operator on (P, \leq) , let $\mathcal{C}_c := \{x \in P \mid x = c(x)\}$.

2 Examples and Exercises

Topological Spaces

Definition 3. *A topological space is a tuple (X, τ) such that X is a set and τ is a collection of subsets of X , called open sets, such that:*

- (a) $\emptyset, X \in \tau$;
- (b) τ is closed under finite intersection and arbitrary unions.

Exercise 1. *Prove that, for every topological space (X, τ) , the map $c : \mathcal{P}(X) \rightarrow \mathcal{P}(X)$ defined by the assignment $Y \mapsto \bigcap \{C \in \mathcal{C}_\tau \mid Y \subseteq C\}$ is a closure operator on $(\mathcal{P}(X), \subseteq)$, where $\mathcal{C}_\tau := \{A^c \mid A \in \tau\}$ and A^c is the relative complement of A with respect to X .*

Logical Systems

Definition 4. The Polish school in logic defines a logical system as a tuple $\mathcal{S} = (\mathbf{Fm}, \vdash)$ such that \mathbf{Fm} is the algebra of \mathcal{L} -formulas over a given set Φ of propositional variables for a given algebraic signature \mathcal{L} , and \vdash is a consequence relation on \mathbf{Fm} , i.e. $\vdash \subseteq \mathcal{P}(\mathbf{Fm}) \times \mathbf{Fm}$ such that for all $\varphi, \psi \in \mathbf{Fm}$ and all $\Delta, \Gamma \subseteq \mathbf{Fm}$:

- (a) if $\varphi \in \Gamma$, then $\Gamma \vdash \varphi$;
- (b) if $\Gamma \vdash \varphi$ and $\Gamma \subseteq \Delta$, then $\Delta \vdash \varphi$;
- (c) if $\Gamma \vdash \psi$ for every $\psi \in \Delta$ and $\Delta \vdash \varphi$, then $\Gamma \vdash \varphi$.

Exercise 2. For any logical system \mathcal{S} as above, prove that the map $c_{\vdash} : \mathcal{P}(\mathbf{Fm}) \rightarrow \mathcal{P}(\mathbf{Fm})$ defined by the assignment $\Gamma \mapsto \{\varphi \mid \Gamma \vdash \varphi\}$ is a closure operator on $(\mathcal{P}(\mathbf{Fm}), \subseteq)$.

Galois Connections

Definition 5 (Galois connection). Let (P, \leq) and (Q, \preceq) be partial orders. A Galois connection is a pair of maps $\triangleright : P \rightarrow Q$ and $\blacktriangleright : Q \rightarrow P$ such that, for every $x \in P$ and every $y \in Q$,

$$x \leq \triangleright y \quad \text{iff} \quad y \preceq \blacktriangleright x.$$

Exercise 3. Prove that, in any Galois connection (cf. Definition ??),

1. $x \leq \triangleright \blacktriangleright x$ and $y \leq \blacktriangleright \triangleright y$;
2. $x \leq x'$ implies $\triangleright x' \preceq \triangleright x$, and $y \preceq y'$ implies $\blacktriangleright y' \leq \blacktriangleright y$;
3. $\triangleright \blacktriangleright \triangleright x = \triangleright x$ and $\blacktriangleright \triangleright \blacktriangleright y = \blacktriangleright y$.

Deduce from the previous items that for every Galois connection as above, $\triangleright \blacktriangleright$ is a closure operator on (P, \leq) and $\blacktriangleright \triangleright$ is a closure operator on (Q, \preceq) .

Polarities

Definition 6. A polarity or formal context is a triple $\mathbb{P} = (A, X, I)$ such that A and X are sets and $I \subseteq A \times X$. Every polarity induces the pair of maps

$$(\cdot)^\uparrow : \mathcal{P}(A) \rightarrow \mathcal{P}(X) \quad \text{and} \quad (\cdot)^\downarrow : \mathcal{P}(X) \rightarrow \mathcal{P}(A),$$

respectively defined by the assignments

$$B^\uparrow := \{x \in X \mid \forall a (a \in B \Rightarrow aIx)\} \quad \text{and} \quad Y^\downarrow := \{a \in A \mid \forall x (x \in Y \Rightarrow aIx)\}.$$

Exercise 4. Show that

1. the map $\uparrow : \mathcal{P}(A) \rightarrow \mathcal{P}(X)$ and the map $\downarrow : \mathcal{P}(X) \rightarrow \mathcal{P}(A)$ form a Galois connection, that is, for every $B \subseteq A$ and every $Y \subseteq X$,

$$B \subseteq Y^\downarrow \quad \text{iff} \quad Y \subseteq B^\uparrow$$

2. deduce that $(\cdot)^{\uparrow\downarrow} : \mathcal{P}(A) \rightarrow \mathcal{P}(A)$ and $(\cdot)^{\downarrow\uparrow} : \mathcal{P}(X) \rightarrow \mathcal{P}(X)$ are closure operators on $(\mathcal{P}(A), \subseteq)$ and $(\mathcal{P}(X), \subseteq)$ respectively.¹

¹When $B = \{a\}$ (resp. $Y = \{x\}$) we write $a^{\uparrow\downarrow}$ for $\{a\}^{\uparrow\downarrow}$ (resp. $x^{\downarrow\uparrow}$ for $\{x\}^{\downarrow\uparrow}$).